

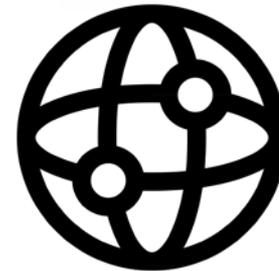
# fling: A Flexible Ping for Middlebox Measurements

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# Motivation

- Lack of a generic tool that can assess whether an arbitrary communication pattern between end points would succeed

Will my new protocol/protocol-extension be blocked or modified by middleboxes?

# fling (flexible-Ping) is an end-to-end active measurement tool

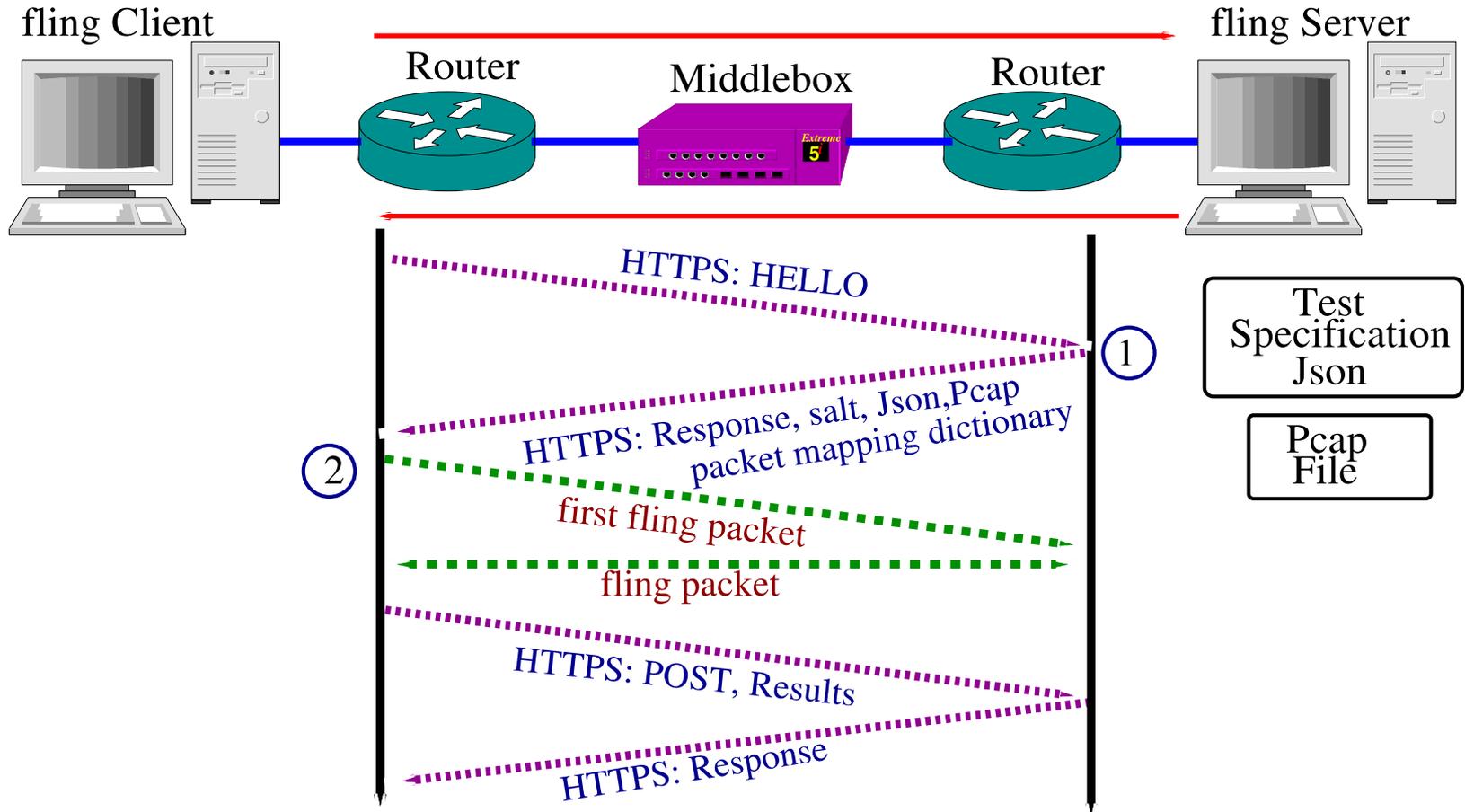
- Allows testing whether an arbitrary sequence of packets can be exchanged between a fling client and a fling server
- Uses raw sockets and supports both IPv4 and IPv6
- Tests needs to be only specified at the server side
- Can narrow down the location of packet modification or drop

# Middleboxes measurement tools

<i>Tool</i>	Raw sockets	Test protocols other than TCP	Test update: need to change	Fully controlled client-server dialogue	Detecting location of modification and drop
<i>fling</i>	✓	✓	Server	✓	✓
<b>Netalyzr</b>	✗	✓*	Server	✓	✓‡
<b>TCPExposure</b>	✓	✗	Both	✓	✗
<b>HICCUPS</b>	✓	✗	Both	✓	✗
<b>Tracebox</b>	✓	✓	Client	✗	✓‡
<b>PATHspider</b>	✓	✓	Client	✗	✓‡
<b>TBit</b>	✓	✗	Client	✗	✗

**Table:** Comparison of related tools. \*ICMP,UDP; ‡One-sided only

# How does it work?

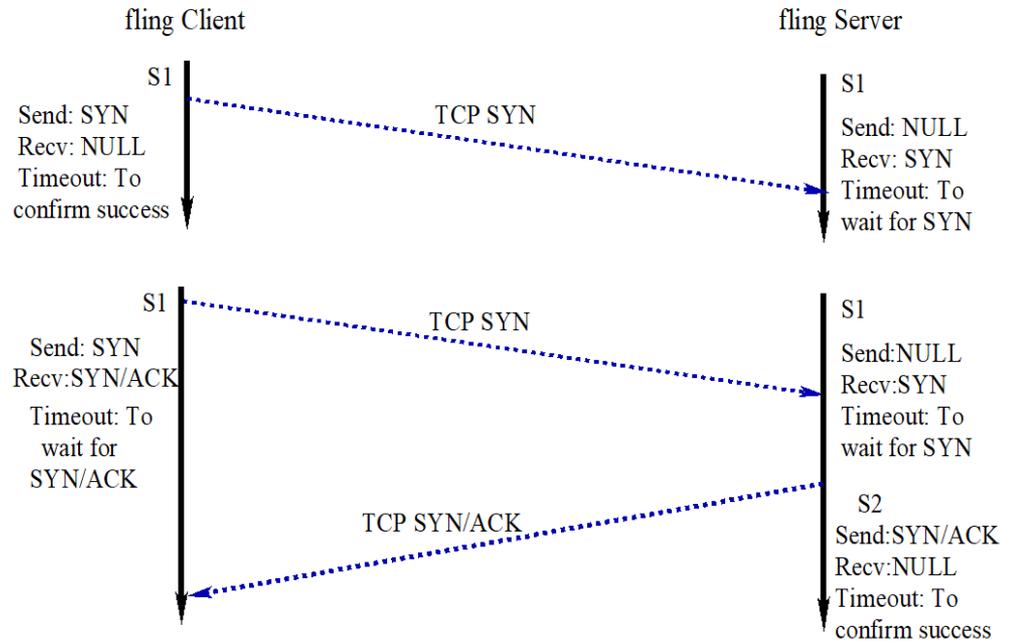


# How does it work?

```
{
  "name": "TCP SYN/ACK test",
  "__index": {
    "0": "TCP SYN",
    "1": "TCP SYN/ACK"
  },
  "packet_Info": [
    {
      "name": "TCP SYN",
      "swap": [0, 2, 2],
      "ChksumType": "adler-32",
      "ChksumPos": [16, 0],
      "ChksumLen": [2, 0],
      "ChksumPseudoHDR": true
    },
    {
      "name": "TCP SYN/ACK",
      "swap": [0, 2, 2],
      "ChksumType": "adler-32",
      "ChksumPos": [16, 0],
      "ChksumLen": [2, 0],
      "ChksumPseudoHDR": true
    }
  ],
  "client": {
    "state_sequence": ["S1"],
    "states": [
      {
        "state": "S1",
        "send": ["TCP SYN"],
        "recv": ["TCP SYN/ACK"],
        "delaySend": [0],
        "timeout": [2000]
      }
    ]
  },
  "server": {
    "state_sequence": ["S1", "S2"],
    "states": [
      {
        "state": "S1",
        "recv": ["TCP SYN"],
        "timeout": [2000]
      },
      {
        "state": "S2",
        "send": ["TCP SYN/ACK"],
        "delaySend": [0],
        "timeout": [2000]
      }
    ]
  }
}
```

(a) json file for a simple TCP SYN-SYN/ACK dialogue test

Pcap file contains a list of packets and Json files describes a dialogue



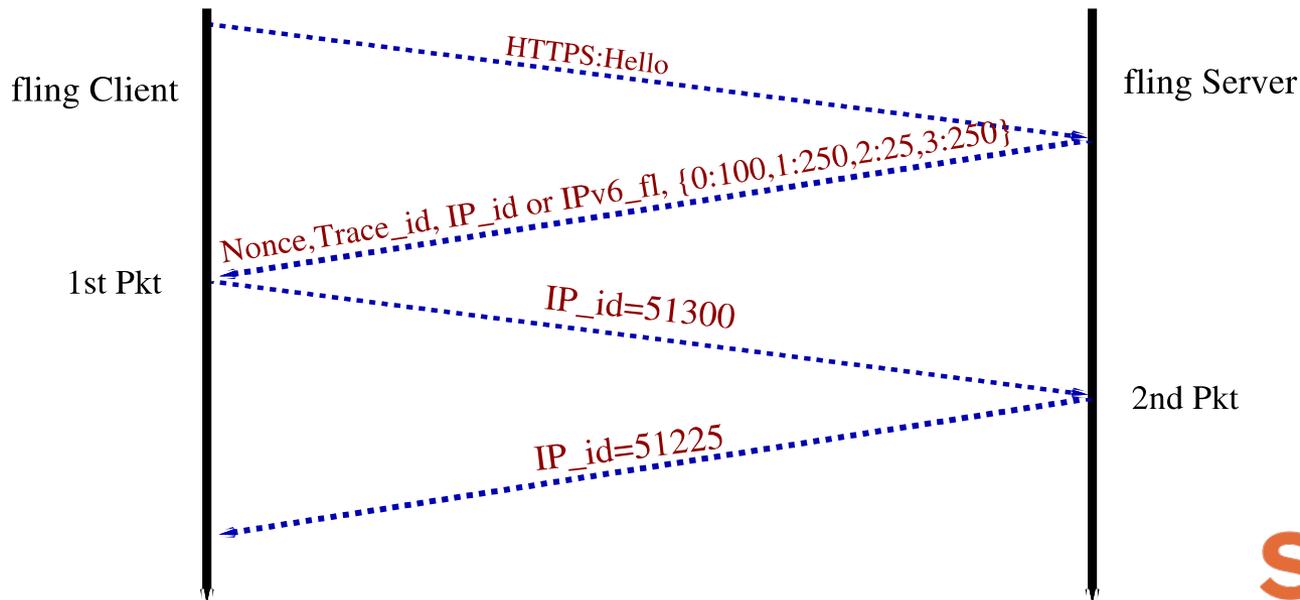
(b) Protocol instance

# Challenges

- Mapping packets into corresponding test sequence
- Detect whether packets are really dropped
- Infer the location of packet modification or drop

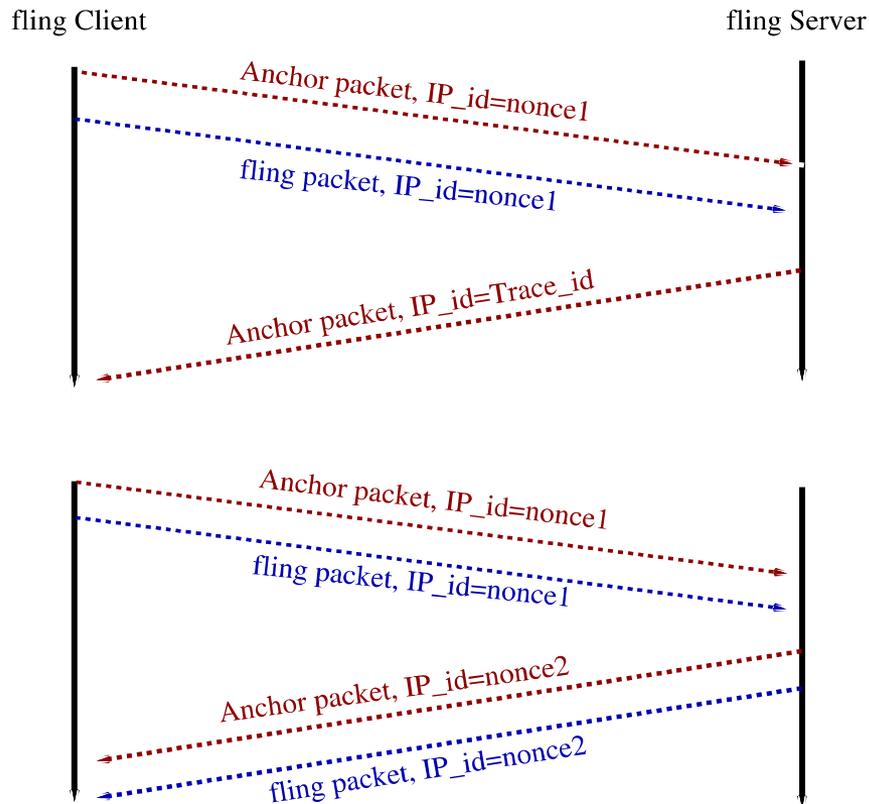
# How to identify packets that belong to a particular test?

- fling uses nonce and protocol numbers for packet identification
- The packet's nonce is (salt,random\_number)
  - Salt is 8-bit number generated by the server for each test
  - The server also generates a random number for each packet
  - The nonce position in the packet is defined in the Json file

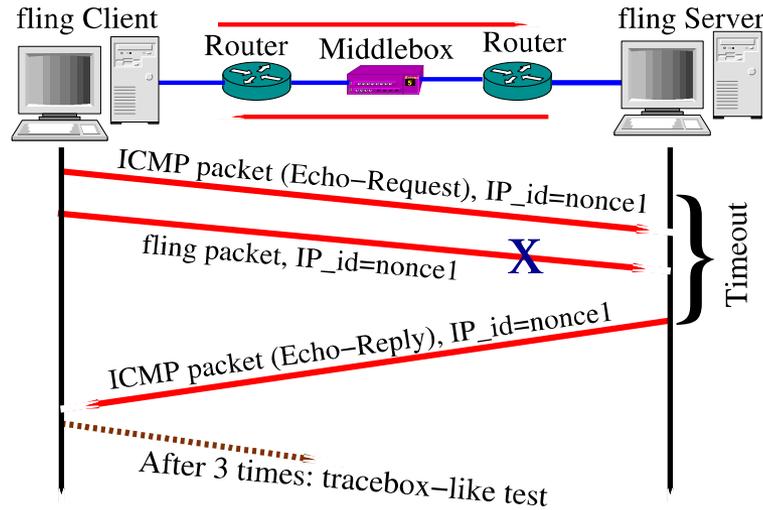


# Detect the drop of test packets

- To confirm drops of tests packets fling sends, along with every test packet, an anchor packet (TCP SYN-SYN/ACK, UDP, or ICMP)



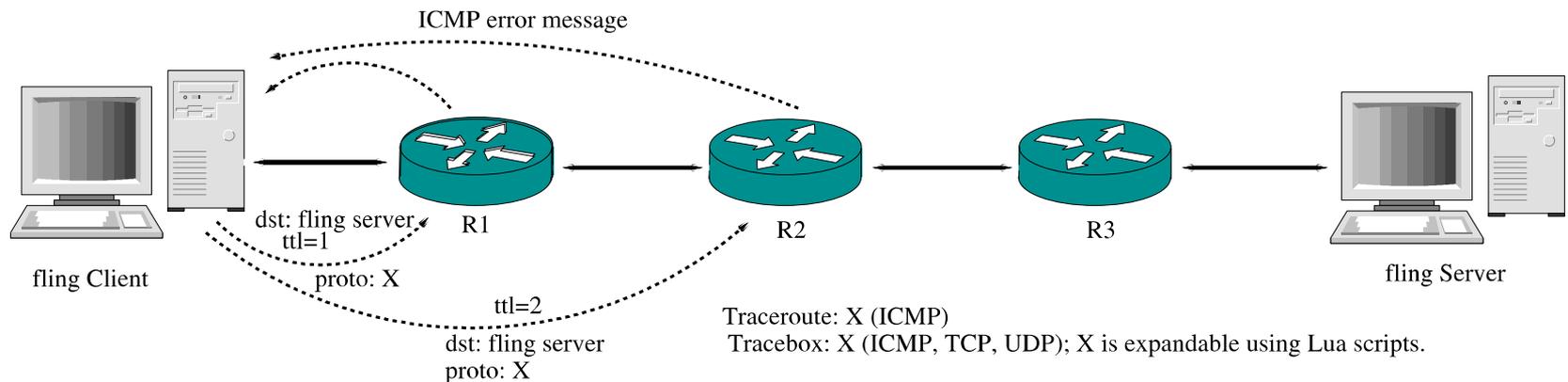
# Detect the drop of test packets



		X DROP; start tracebox-like test
DROPPED	DROPPED	Repeat 3×. Then, assume: CONGESTION; start tracebox-like test

# Detecting the location of modification or drop

- In case a test packet is dropped fling re-sends the test packet with an increasing TTL
- RFC 1812- compliant routers enclose the entire packet in the payload of the ICMP error message



# Case study: uses fling to check whether DSCP code points survive end-to-end paths

- WebRTC would like use DSCP code-points to signal QoS expectations but does it really work?
- We tested three DSCP values: CS1 (low priority data), AF42 (Multimedia conferencing) and EF (Telephony)



## Clients

Testbed	IPv4	IPv6
Ark	111	46
NorNet Core	40	19
PlanetLab	14	-

34 IPv4 servers  
18 IPv6 servers

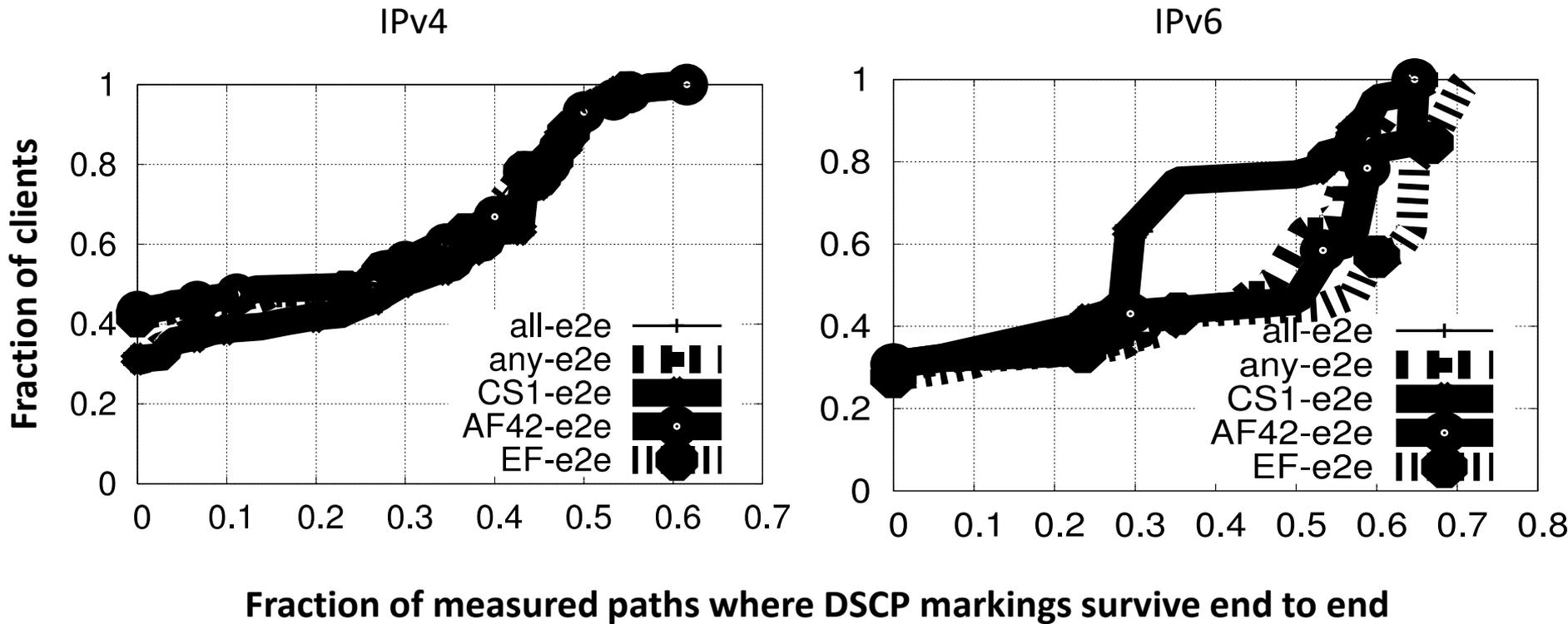
~10K IPv4 paths  
~2K IPv6 paths

298 IPv4 ASES and 119 IPv6 ASES

All key large transit providers + many access providers e.g. ComCast, Bharti AirTel and CenturyLink.

# Case study: uses fling to check whether DSCP code points survive end-to-end paths

DSCP markings survived e2e in 33% and 50% for IPv4 and IPv6, respectively

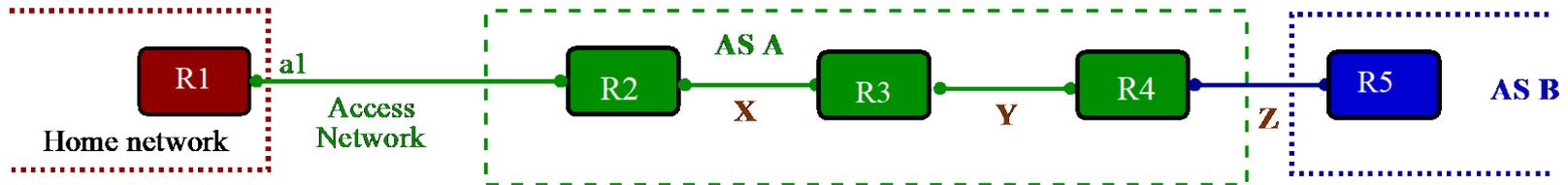


# Do packets with DSCP markings risk being blackholed?

Code Point	Direction	Total failures	# clients	#servers
CS1	Forward	18	6	10
CS1	Reverse	74	27	31
AF42	Forward	28	9	16
AF42	Reverse	74	27	28
EF	Forward	28	9	17
EF	Reverse	76	23	32
All	Forward	13	3	6
All	Reverse	27	11	15

None of these failures happened at TTL 1 or 2

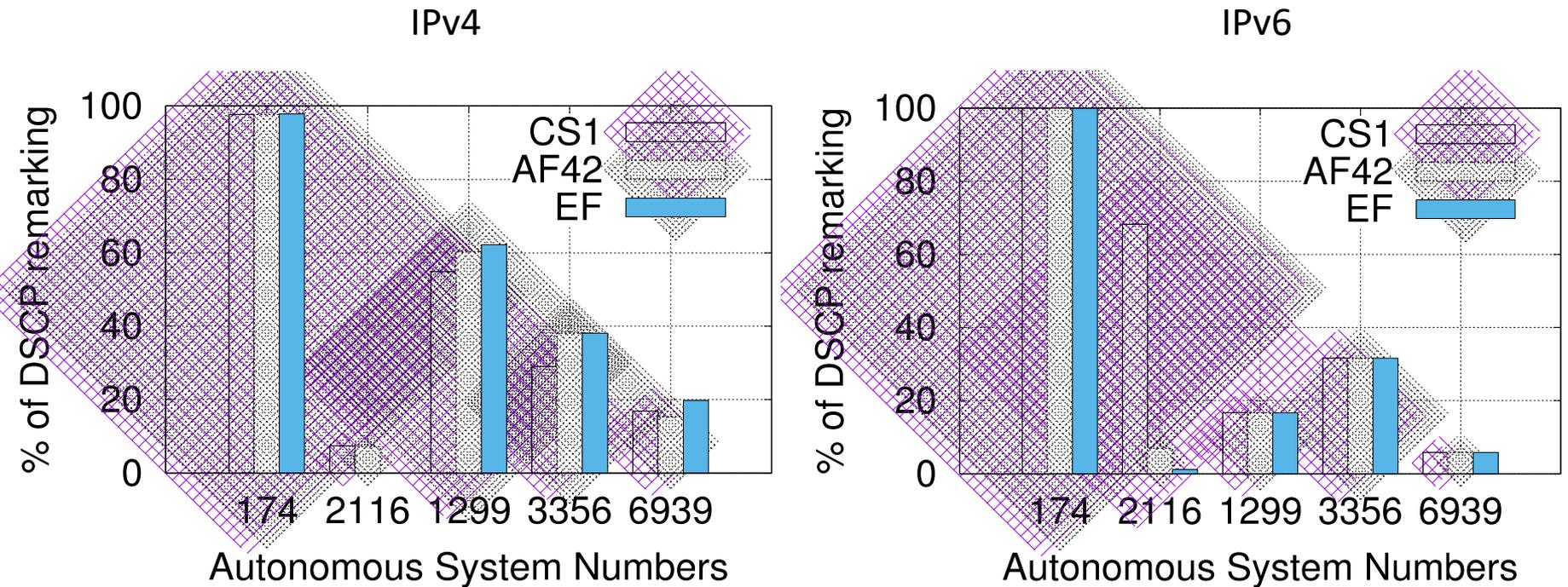
# Where was DSCP re-marked?



Changed in	IPv4	IPv6
Home network	21%	12%
First-hop AS	43%	31%
Beyond the first-hop AS	36%	57%

- Home gateways treats DSCP in a myriad of ways: zero, re-write to unused value, re-write to a used value
- First hop Ases often zero DSCP

# ASes beyond the first-hop AS employ a diverse set of re-marking policies



- Cogent remarks everything to either AF11 or AF21
- Other large ISPs do not seem to modify DSCP markings

# Limitations of the DSCP study

- Although we have around 10k paths, the coverage remains sparse
- The fact that DSCP marking survives does not imply that marked traffic will be treated differently
- All probes are in fixed networks

# Takeaways

- fling is a flexible tool that allows for a wide range of middlebox tests
- We have used fling to investigate whether DSCP markings survive routers and middleboxes
- Please help us increasing our coverage by running fling (email me [ahmed@simula.no](mailto:ahmed@simula.no))