Yarrp'ing the Internet

Robert Beverly

Naval Postgraduate School

February 12, 2016

Active Internet Measurements (AIMS) Workshop



Active Topology Probing

• Years (and years) of prior work on Internet-scale topology probing

• e.g., Scamper, DoubleTree, iPlane

lt's 2016:

Why can't we traceroute to every IPv4 destination quickly?

- e.g., O(minutes)?
- (The ZMap^a and Masscan^b folks can do it why can't we?)

^aZ. Durumeric et al., 2013 ^bR. Graham, 2013



Active Topology Probing

• Years (and years) of prior work on Internet-scale topology probing

e.g., Scamper, DoubleTree, iPlane

lt's 2016:

- Why can't we traceroute to every IPv4 destination quickly?
- e.g., O(minutes)?
- (The ZMap^a and Masscan^b folks can do it why can't we?)

^aZ. Durumeric et al., 2013 ^bR. Graham, 2013



Existing traceroute-style approaches:

- Maintain state over outstanding probes (identifier, origination time)
- Are sequential, probing all hops along the path. At best, parallelism limited to a window of outstanding destinations being probed.

Implications:

- **Concentrates load:** along paths, links, routers (potentially triggering rate-limiting or IDS alarms)
- Production systems probe slowly



Existing traceroute-style approaches:

- Maintain state over outstanding probes (identifier, origination time)
- Are sequential, probing all hops along the path. At best, parallelism limited to a window of outstanding destinations being probed.

Implications:

- **Concentrates load:** along paths, links, routers (potentially triggering rate-limiting or IDS alarms)
- Production systems probe slowly





"Yelling at Random Routers Progressively"

Takes inspiration from ZMap:

- Uses a block cipher to randomly permute the < IP, TTL > space
- Is stateless, recovering necessary information from replies
- Permits fast Internet-scale active topology probing (even from a single VP)



T,

13

Example Topology

prober



R. Beverly (NPS)

AIMS 2016 5 / 17

 T_2



Traditional traceroute sends probes with incrementing TTL to destination T_1





... continuing until finished with T_1 (reach destination or gap limit). Prober must maintain state, while traffic is concentrated on *prober* \rightsquigarrow T_1 path



Yarrp iterates through randomly permuted < Target, TTL > pairs





Yarrp iterates through randomly permuted < Target, TTL > pairs



Yarrp

Inferred Topology



Finally, stitch together topology. Requires state and computation, but off-line after probing completes.

n. Develly (INFO	R.	Beverl	v (NPS
------------------	----	--------	-----	------------

Methodology

Challenges

Encoding State



- IPID = Probe's TTL
- TCP Source Port = cksum(Target IP destination)^a
- TCP Seq No = Probe send time (elapsed ms)
- Per-flow load balancing fields remain constant (ala Paris)
- Assume routers echo only 28B of expired packet

^aMalone PAM 2007: \approx 2% of quotations contained modified destination IP

Methodology

Challenges

Recovering State



ICMP TTL exceeded replies permit recovery of: target probed, originating TTL (hop), and responding router interface at that hop.

Challenges

Distribution of unique interfaces discovered vs. TTL for all Ark monitors, one Ark topology probing cycle



Problem: knowing when to stop Little discoverable topology past TTL=32 $\bullet \Rightarrow \text{limit}$

< IP, TTL >search space to *TTL* < 32

< 17 ▶

Initial Testing Speed

- C++ implementation w/o tuning
- Linux KVM (1 core, Intel L5640 @ 2.27GHz)
- Achieve 106K pps

Proof-of-concept

- Sent 10M probes in \approx 100 sec
- Discovered 178,453 unique router interfaces
- CPU: 52%



What's Possible

Traceroute to an address in each /24, for TTLs 1-32

$$t = \frac{2^{24} * 2^5}{100 K pps} \simeq 84 min$$

Traceroute to every routed IPv4 destination

$$t=rac{2^{31}*2^5}{100$$
 Kpps} $\simeq 1$ week



AIMS 2016

15/17

Optimizations

- Base Yarrp requires no state
- (Must reconstruct traces, but that's an offline local process)
- If we're willing to maintain some space, we can optimize: Time Memory Trade Off
 - Probe only routed destinations (radix trie BGP RIB)
 - Avoiding repeated re-discovery of prober's local neighborhood (state over small number of interfaces near prober)
- Distribute: only requires communicating block cipher key and offset!



Next Steps

Yarrping the Internet

- Push limits on how fast we can map the entire IPv4 Internet
- Compare discovered topologies from e.g. Ark versus Yarrp

Applications?

- What do two snapshots of the Internet topology separated by an hour reveal?
- Others?

Thanks! – Questions?

https://www.cmand.org