Logging System for Longlifetime Data Validation

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 must be limited

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 maintenance is complicated
 - key rollover
 - publishing re-signed data
- Can we sign data once and leave it alone?
 - post-fact validation
 - validate data with an expired key?

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- Can we have a time machine to go back?

- a logging system may help!

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 - tamper-evident log

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hashk, 0

hash1, 2

. . .

hash02

slot2



hash1, N-2

hash0N

slotN

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- Encode each sub tree in a data packet
 - name: /<loggerPrefix>/[subTreeIndex]/[digest]
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- Leaf node
 - name: /<loggerPrefix>/leaf/[seqNo]
 - detailed info (signed data, timestamp...)



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 - no need to retrieve log for unavailable data
- User cannot change actual data
 digest is fixed in log
- Users may even keep a sub-tree
 - contain a user's own data
 - could be incomplete
 - root digest is fixed in log



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- Loggers synchronize with each other
 - improve redundancy
 - automatically audit each other
 - using/extending ChronoSync
 - each logger has its own prefix & seqNo

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- Temper-evident log is based on hash function
- A hash function may be broken eventually
- Two copies using different hash functions
 one is relatively stronger than the other
 - e.g., Sha256(B), Sha3-384(B)
 - assume: not broken on the same day
 - weaker broken, stronger still valid
 - enough time to reconstruct another copy with a stronger hash at that time
 - hopefully, it rarely happens

Conclusion

- logging system enables
 - post-fact validation
 - usage of short-lived keys
- Secure logging system through public auditing
- Increase redundancy of certificate provisioning

Thank you!

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